

# Max Flow and Min Cut

# Maximum Flow and Minimum Cut

## Max flow and min cut.

- Two very rich algorithmic problems.
- Cornerstone problems in combinatorial optimization.
- Beautiful mathematical duality.

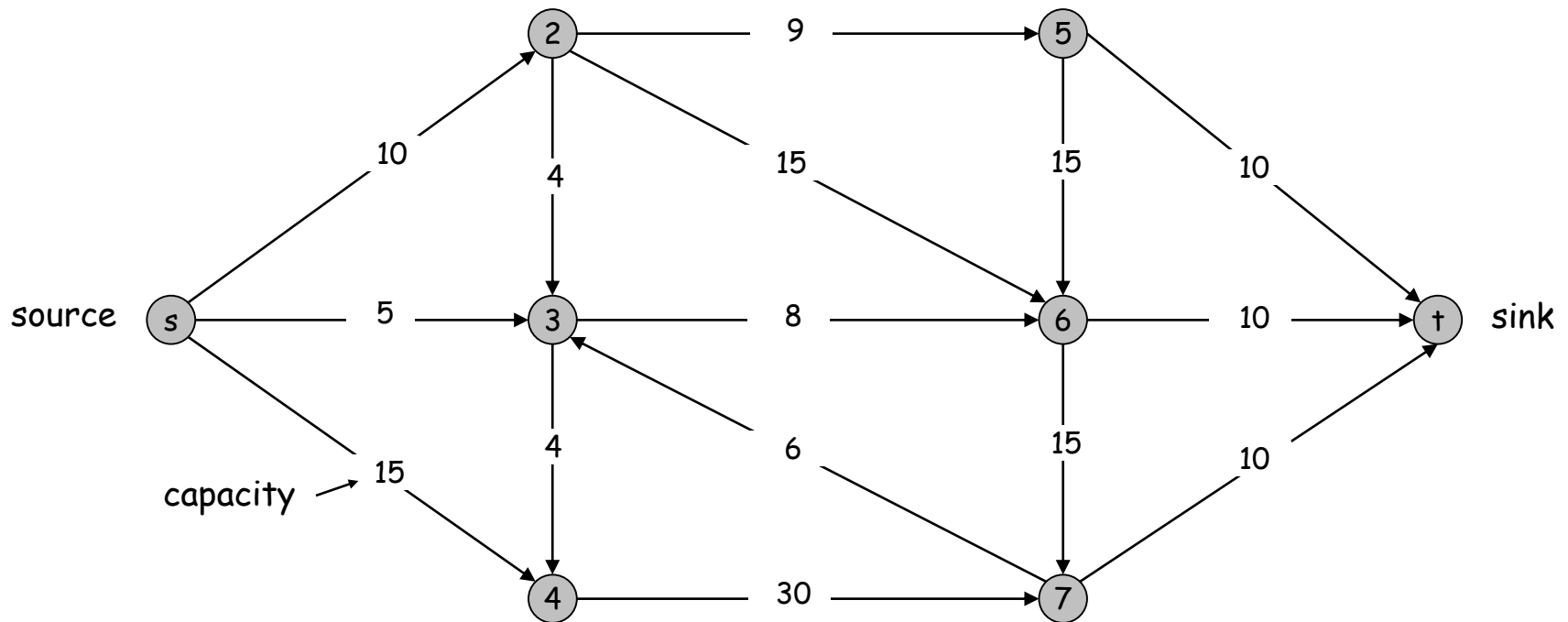
## Nontrivial applications / reductions.

- Data mining.
- Open-pit mining.
- Project selection.
- Airline scheduling.
- Bipartite matching.
- Baseball elimination.
- Image segmentation.
- Network connectivity.
- Network reliability.
- Distributed computing.
- Egalitarian stable matching.
- Security of statistical data.
- Network intrusion detection.
- Multi-camera scene reconstruction.
- Many many more ...

# Minimum Cut Problem

## Flow network.

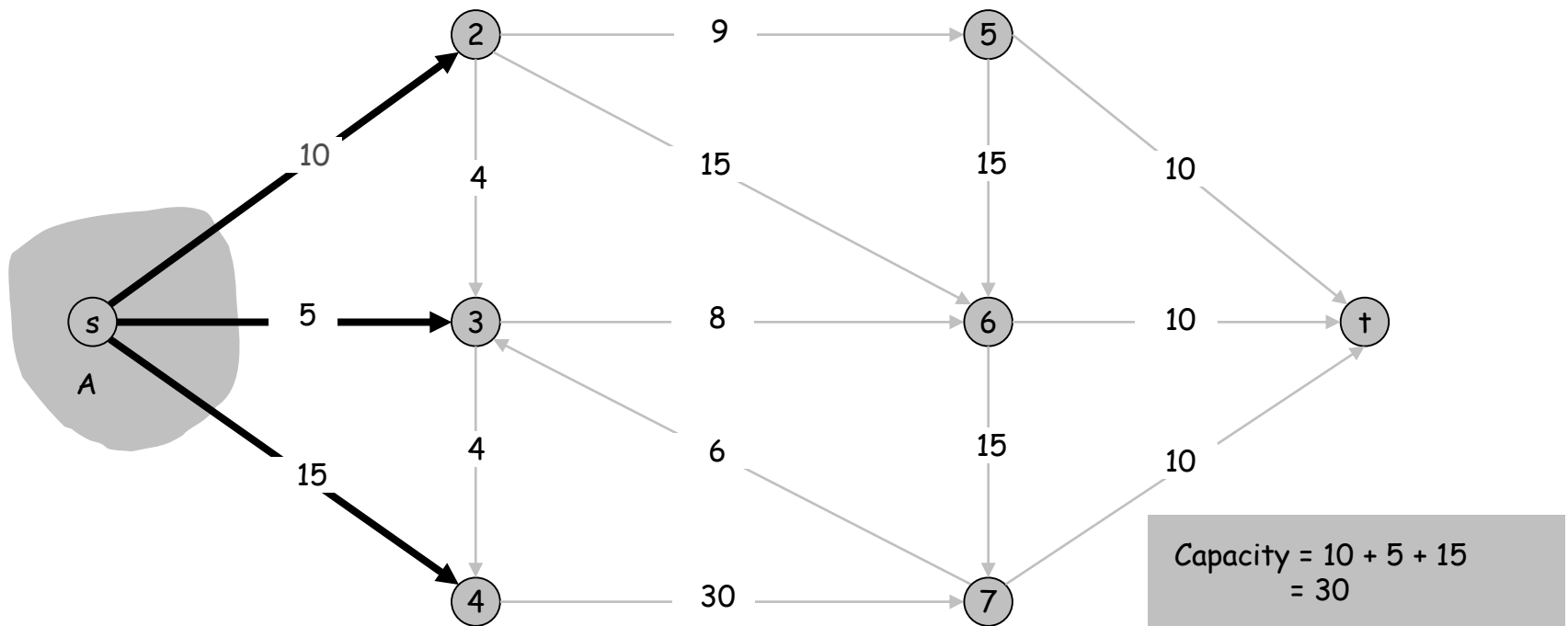
- Abstraction for material **flowing** through the edges.
- $G = (V, E)$  = directed graph, no parallel edges.
- Two distinguished nodes:  $s$  = source,  $t$  = sink.
- $c(e)$  = capacity of edge  $e$ .



# Cuts

Def. An **s-t cut** is a partition  $(A, B)$  of  $V$  with  $s \in A$  and  $t \in B$ .

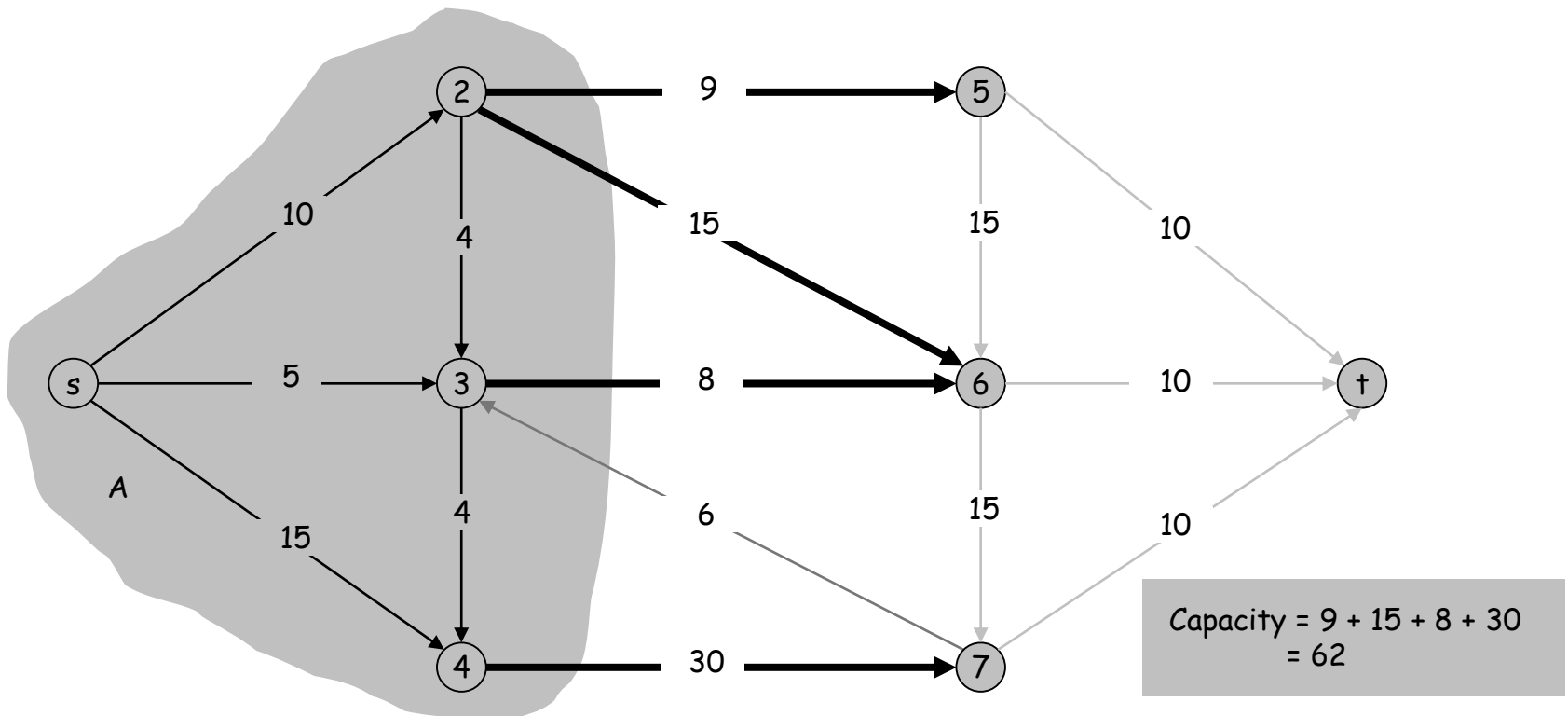
Def. The **capacity** of a cut  $(A, B)$  is:  $cap(A, B) = \sum_{e \text{ out of } A} c(e)$



# Cuts

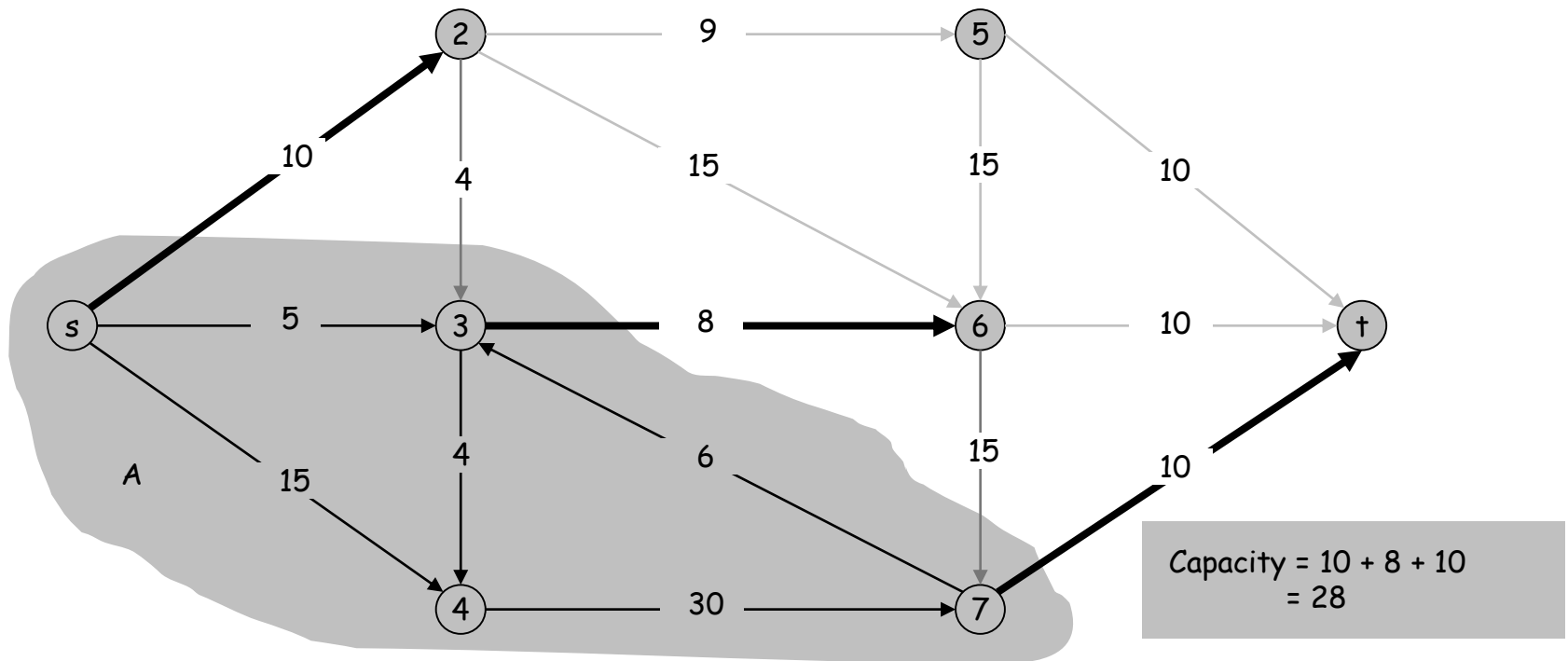
Def. An **s-t cut** is a partition  $(A, B)$  of  $V$  with  $s \in A$  and  $t \in B$ .

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# Minimum Cut Problem

Min s-t cut problem. Find an s-t cut of minimum capacity.

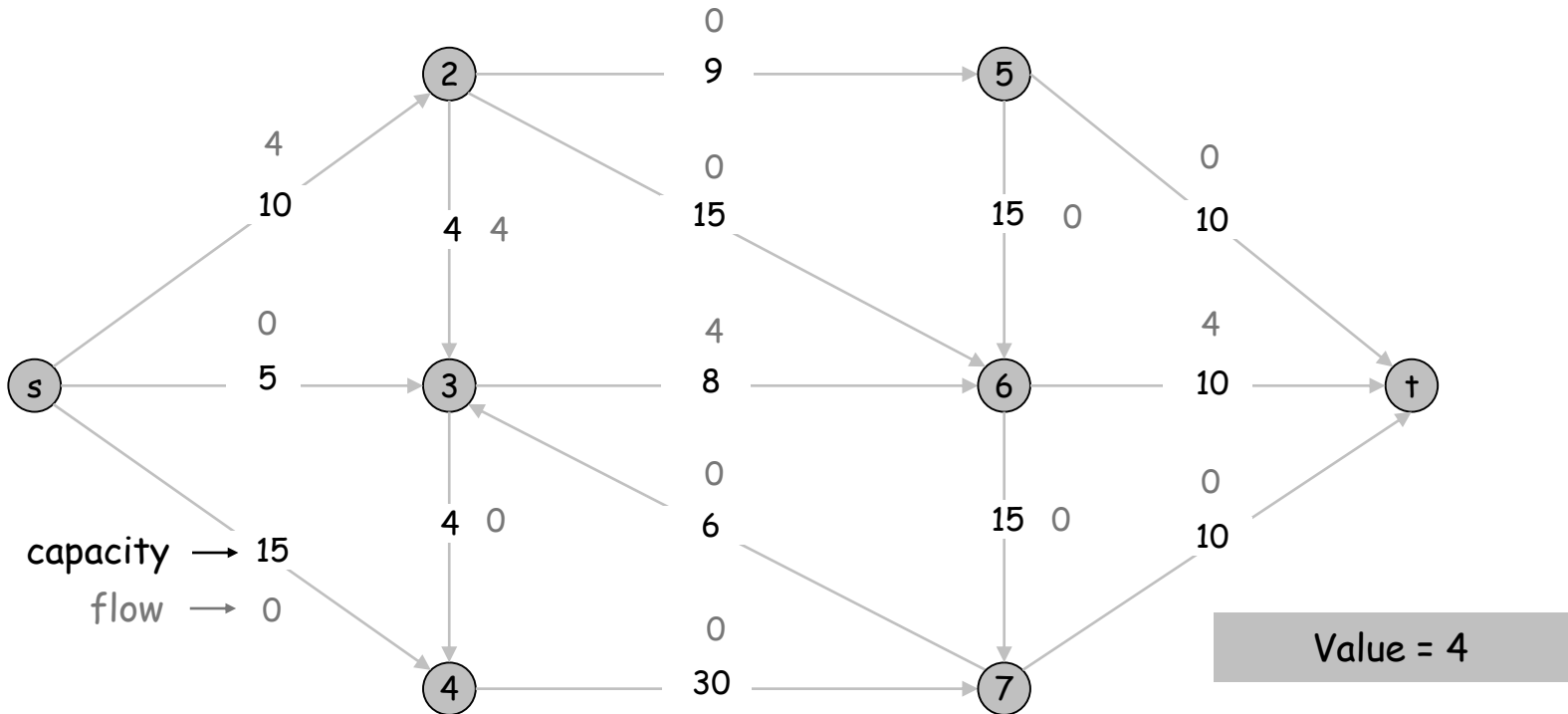


# Flows

Def. An **s-t flow** is a function that satisfies:

- For each  $e \in E$ :  $0 \leq f(e) \leq c(e)$  [capacity]
- For each  $v \in V - \{s, t\}$ :  $\sum_{e \text{ in to } v} f(e) = \sum_{e \text{ out of } v} f(e)$  [conservation]

Def. The **value** of a flow  $f$  is:  $v(f) = \sum_{e \text{ out of } s} f(e)$ .

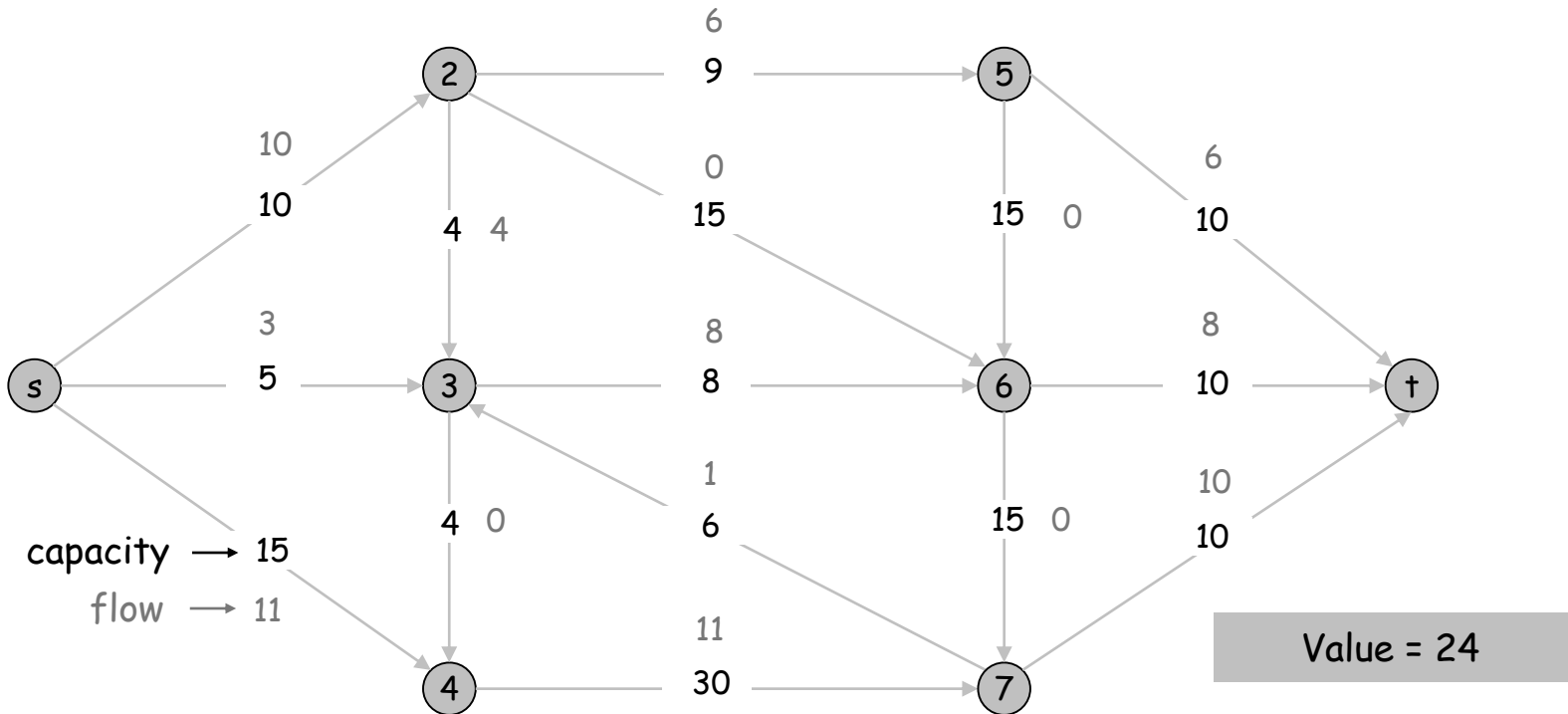


# Flows

Def. An **s-t flow** is a function that satisfies:

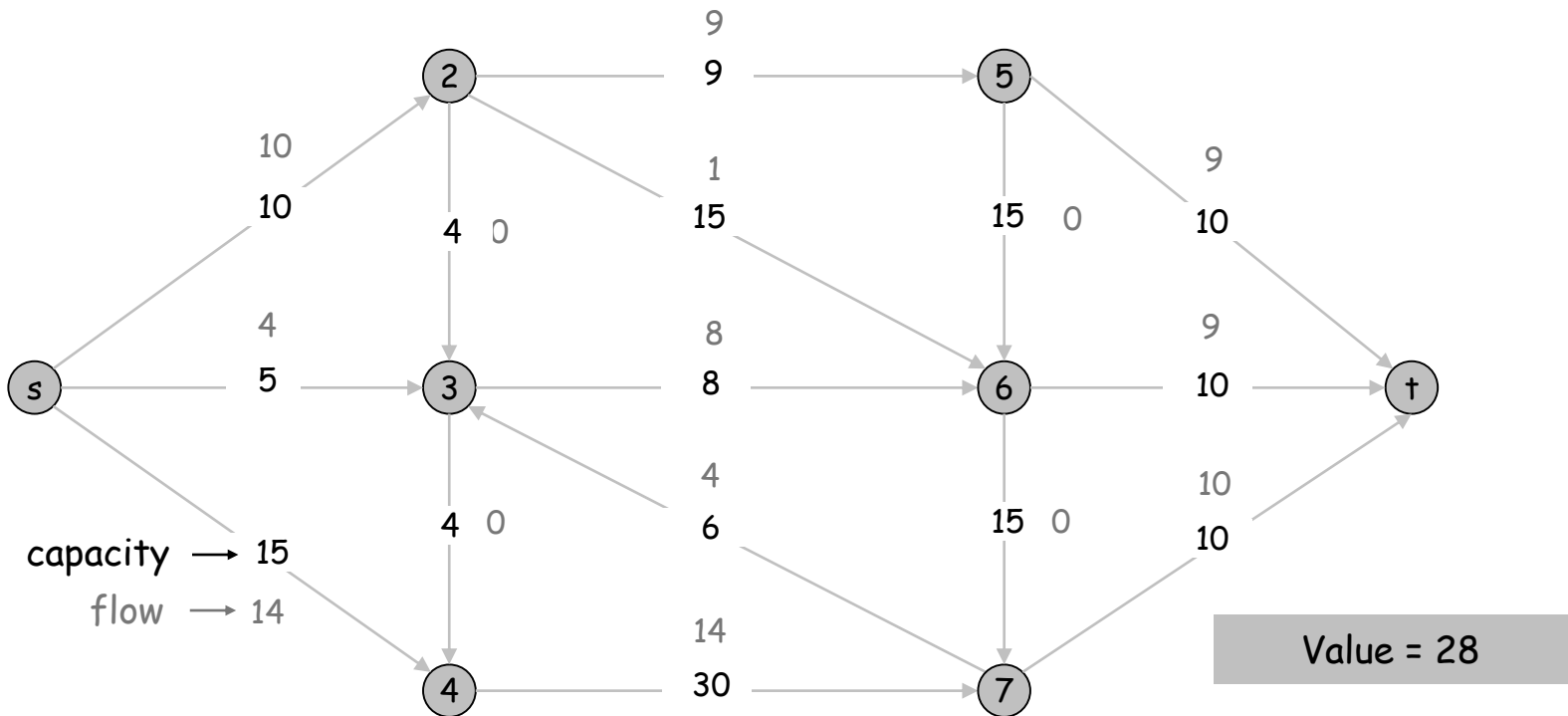
- For each  $e \in E$ :  $0 \leq f(e) \leq c(e)$  [capacity]
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# Maximum Flow Problem

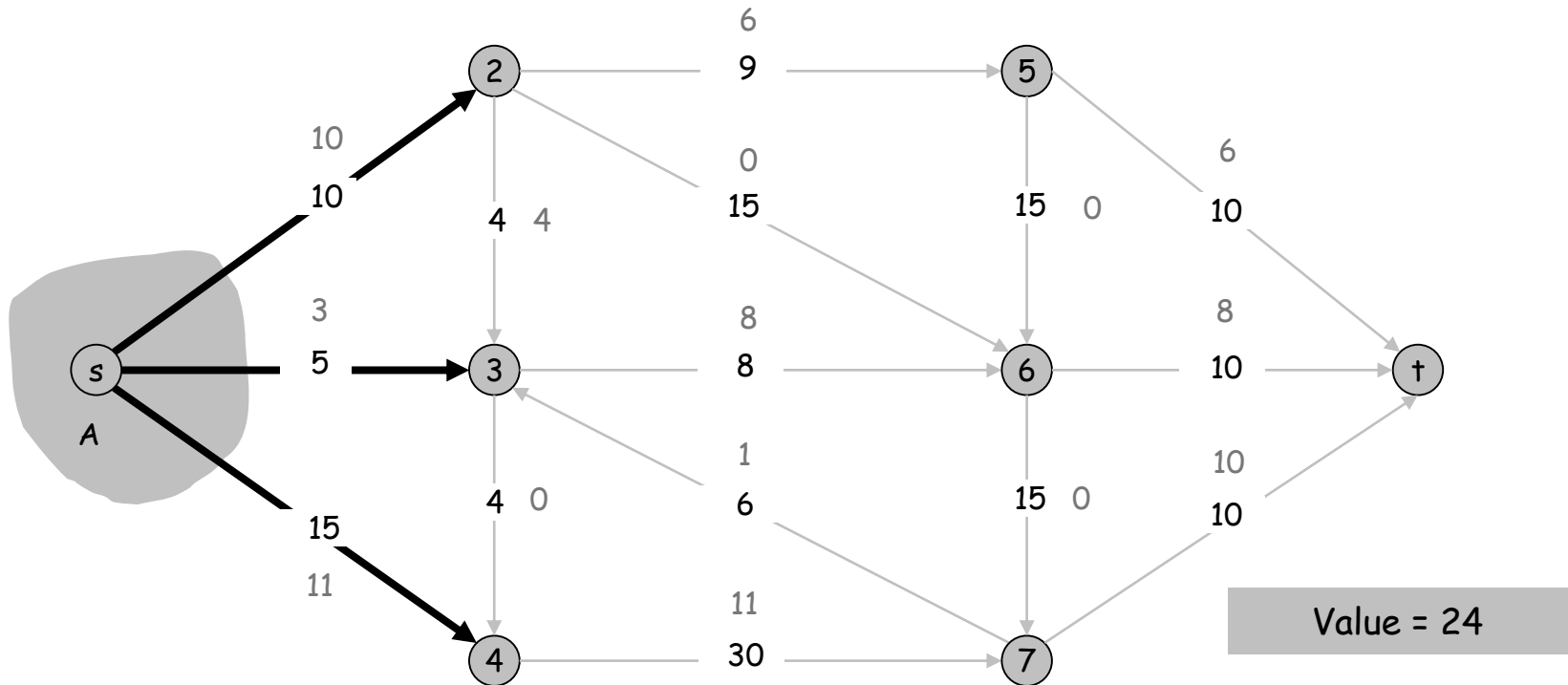
Max flow problem. Find s-t flow of maximum value.



# Flows and Cuts

**Flow value lemma.** Let  $f$  be any flow, and let  $(A, B)$  be any  $s$ - $t$  cut. Then, the net flow sent across the cut is equal to the amount leaving  $s$ .

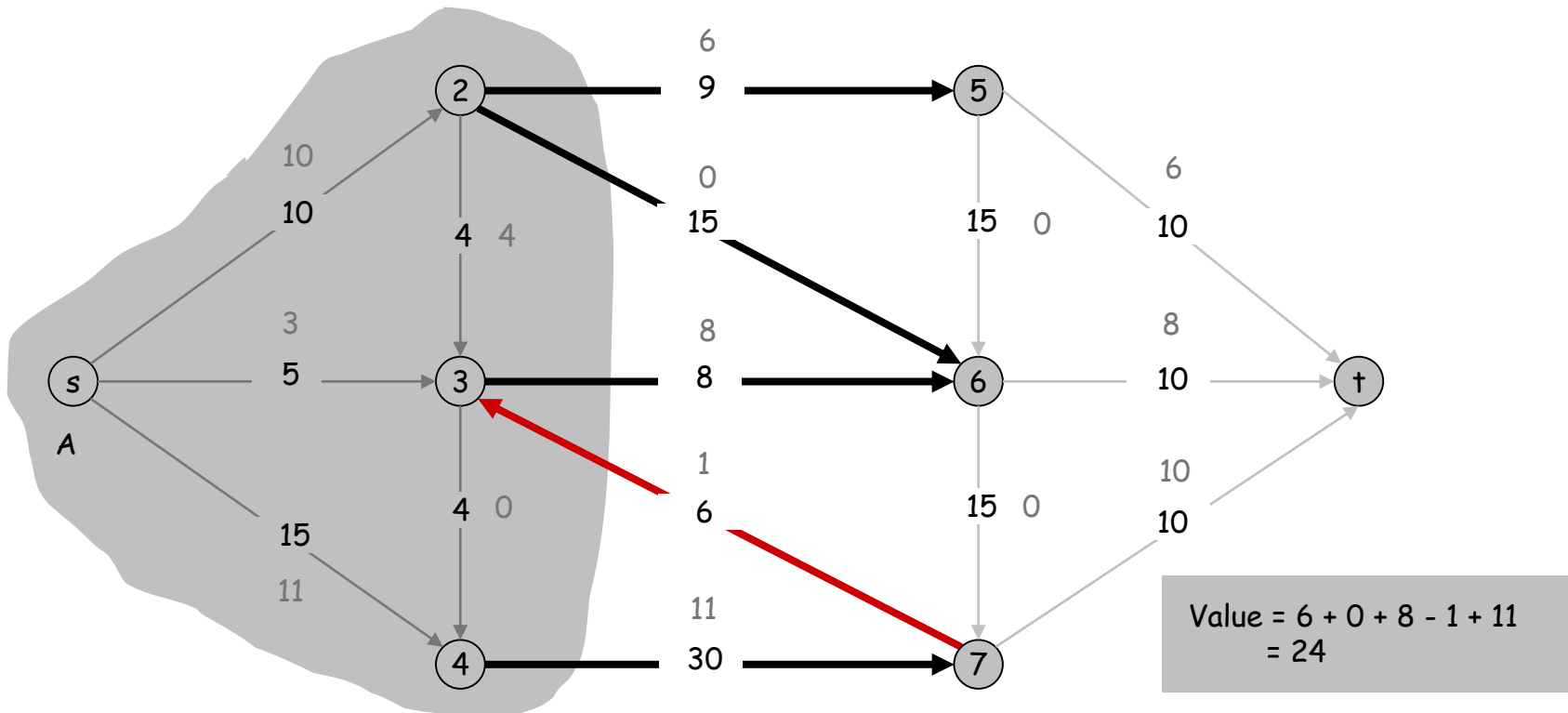
$$\sum_{e \text{ out of } A} f(e) - \sum_{e \text{ in to } A} f(e) = v(f)$$



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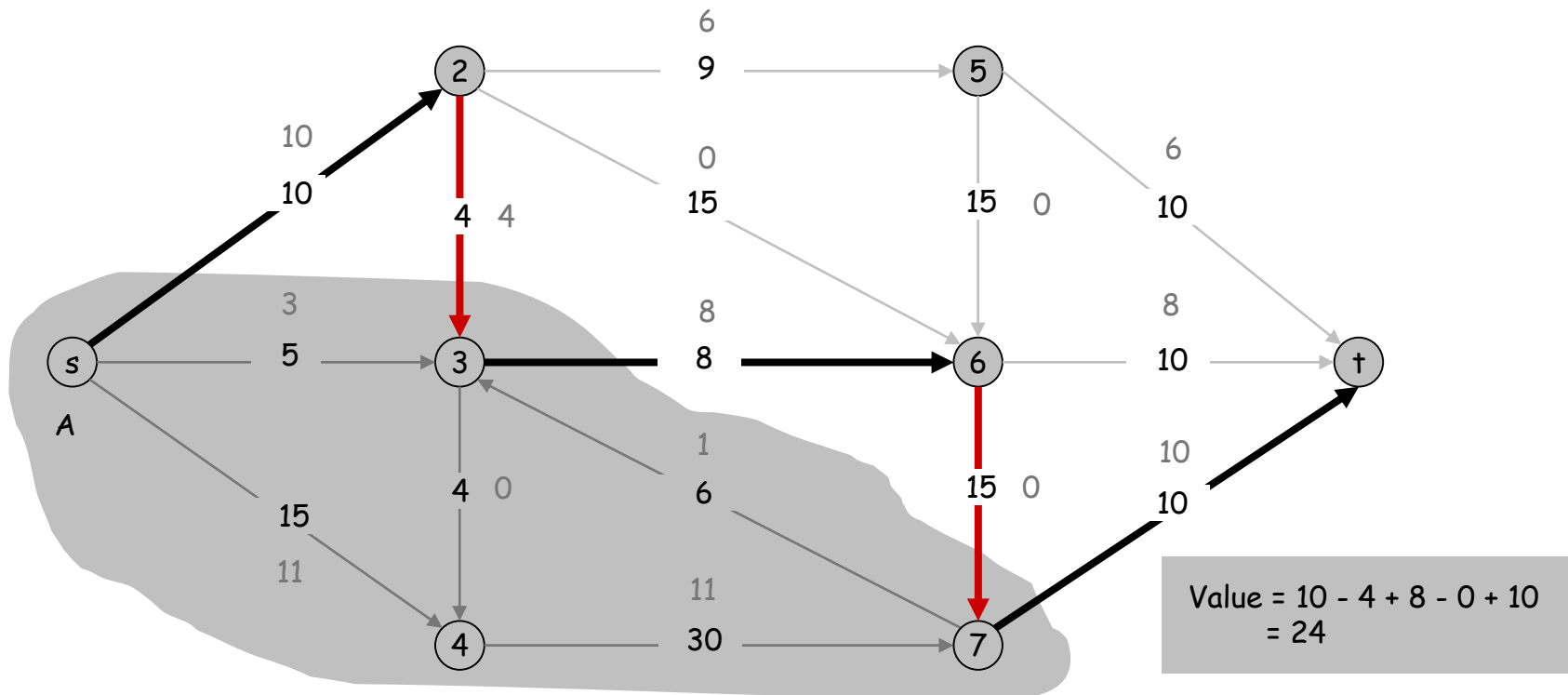
$$\sum_{e \text{ out of } A} f(e) - \sum_{e \text{ in to } A} f(e) = v(f)$$



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# Flows and Cuts

**Flow value lemma.** Let  $f$  be any flow, and let  $(A, B)$  be any  $s$ - $t$  cut. Then

$$\sum_{e \text{ out of } A} f(e) - \sum_{e \text{ in to } A} f(e) = v(f).$$

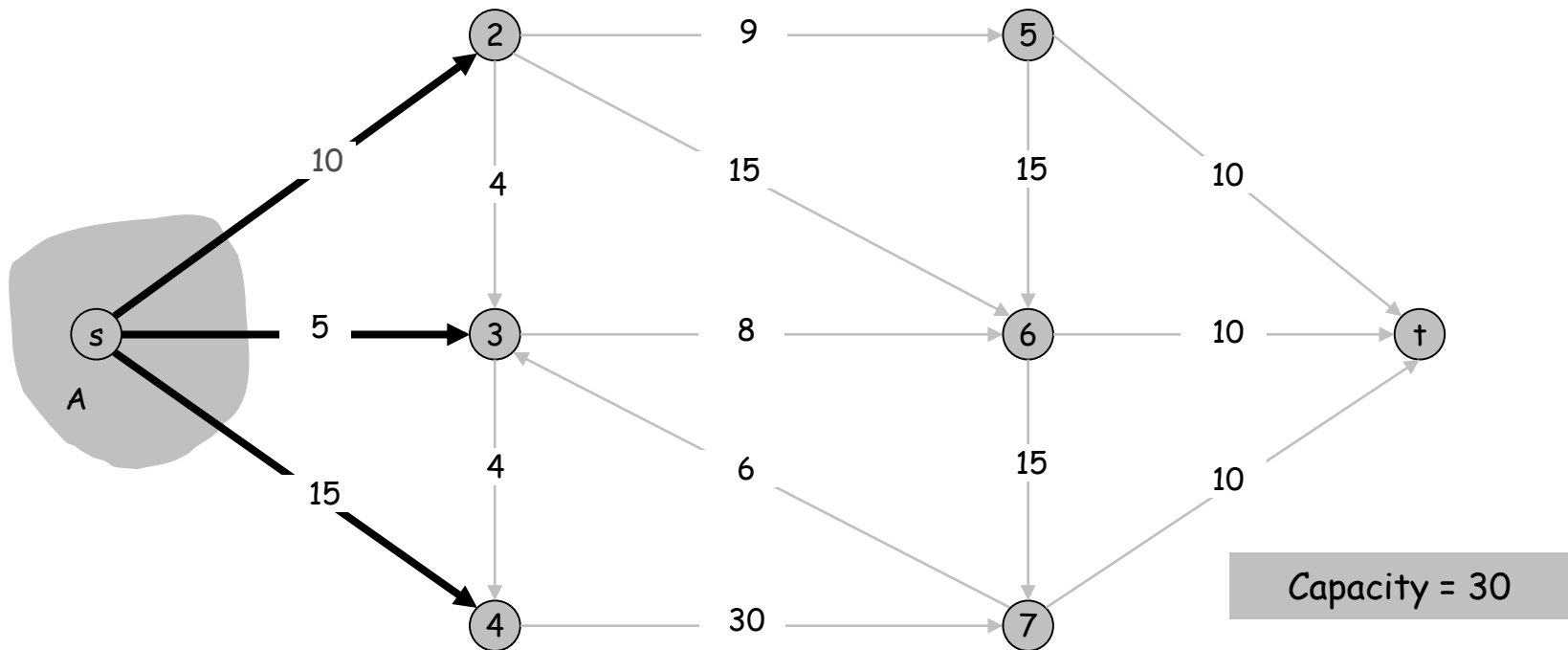
**Pf.**

$$\begin{aligned} v(f) &= \sum_{e \text{ out of } s} f(e) \\ \text{by flow conservation, all terms} &\rightarrow = \sum_{v \in A} \left( \sum_{e \text{ out of } v} f(e) - \sum_{e \text{ in to } v} f(e) \right) \\ \text{except } v = s \text{ are } 0 & \\ &= \sum_{e \text{ out of } A} f(e) - \sum_{e \text{ in to } A} f(e). \end{aligned}$$

# Flows and Cuts

**Weak duality.** Let  $f$  be any flow, and let  $(A, B)$  be any  $s$ - $t$  cut. Then the value of the flow is at most the capacity of the cut.

Cut capacity = 30  $\Rightarrow$  Flow value  $\leq$  30

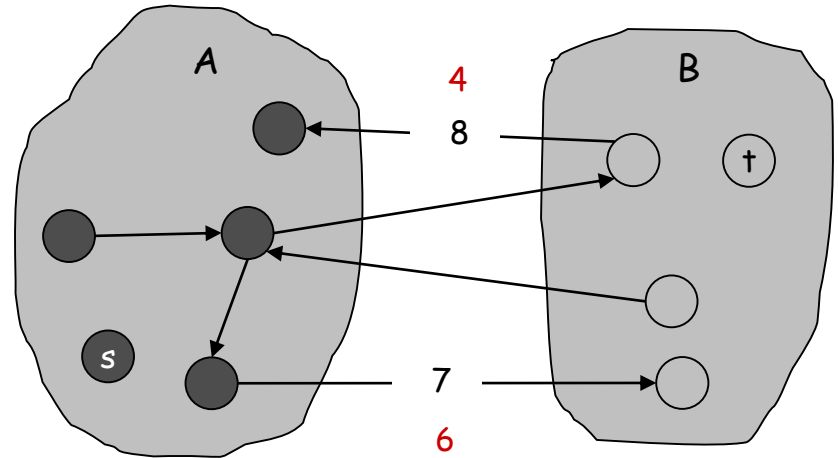


# Flows and Cuts

**Weak duality.** Let  $f$  be any flow. Then, for any  $s$ - $t$  cut  $(A, B)$  we have  $v(f) \leq \text{cap}(A, B)$ .

**Pf.**

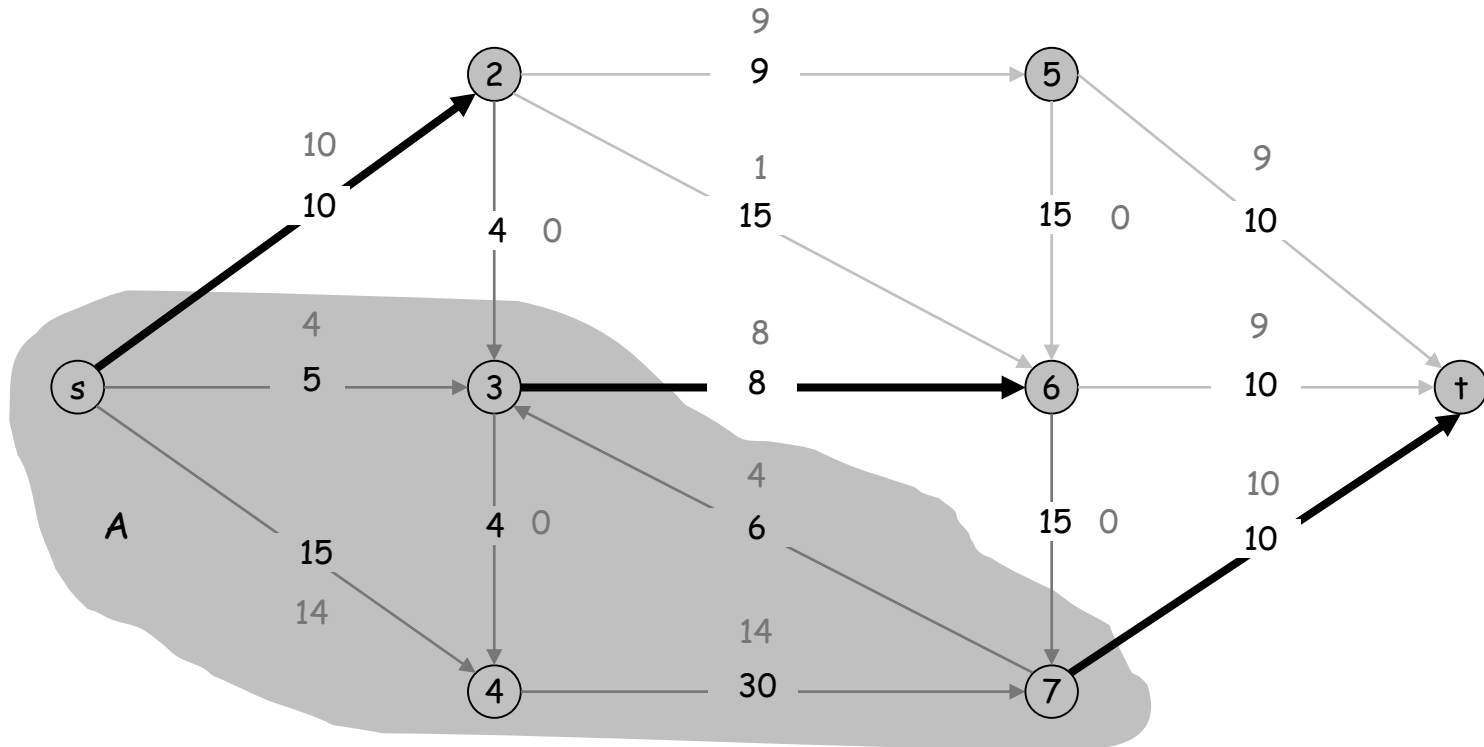
$$\begin{aligned} v(f) &= \sum_{e \text{ out of } A} f(e) - \sum_{e \text{ in to } A} f(e) \\ &\leq \sum_{e \text{ out of } A} c(e) \\ &\leq \sum_{e \text{ out of } A} c(e) \\ &= \text{cap}(A, B) \quad \blacksquare \end{aligned}$$



# Certificate of Optimality

**Corollary.** Let  $f$  be any flow, and let  $(A, B)$  be any cut. If  $v(f) = \text{cap}(A, B)$ , then  $f$  is a max flow and  $(A, B)$  is a min cut.

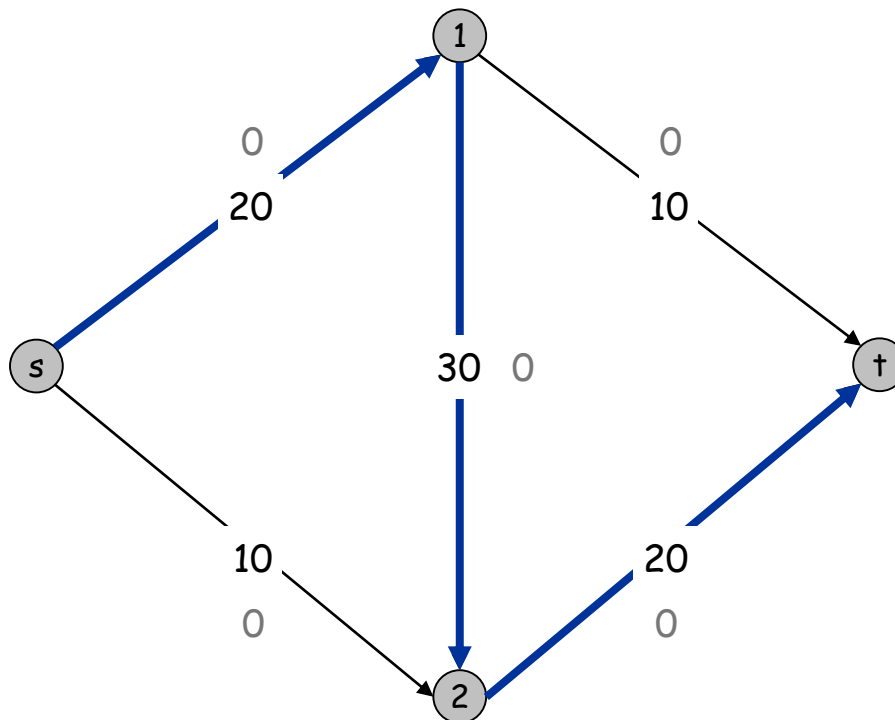
Value of flow = 28  
 Cut capacity = 28  $\Rightarrow$  Flow value  $\leq$  28



# Towards a Max Flow Algorithm

## Greedy algorithm.

- Start with  $f(e) = 0$  for all edge  $e \in E$ .
- Find an  $s$ - $t$  path  $P$  where each edge has  $f(e) < c(e)$ .
- Augment flow along path  $P$ .
- Repeat until you get stuck.

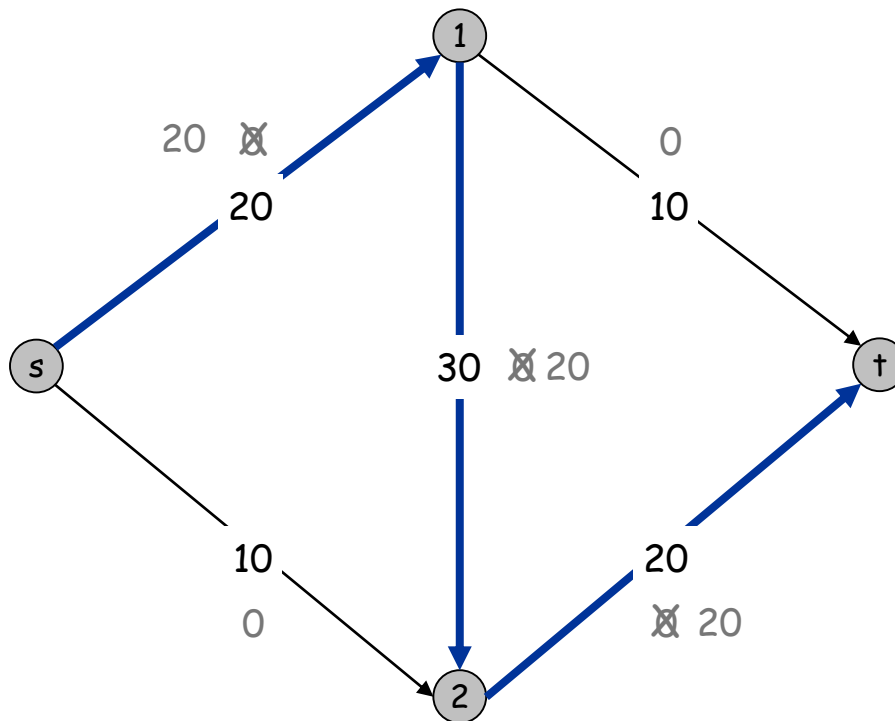


Flow value = 0

# Towards a Max Flow Algorithm

## Greedy algorithm.

- Start with  $f(e) = 0$  for all edge  $e \in E$ .
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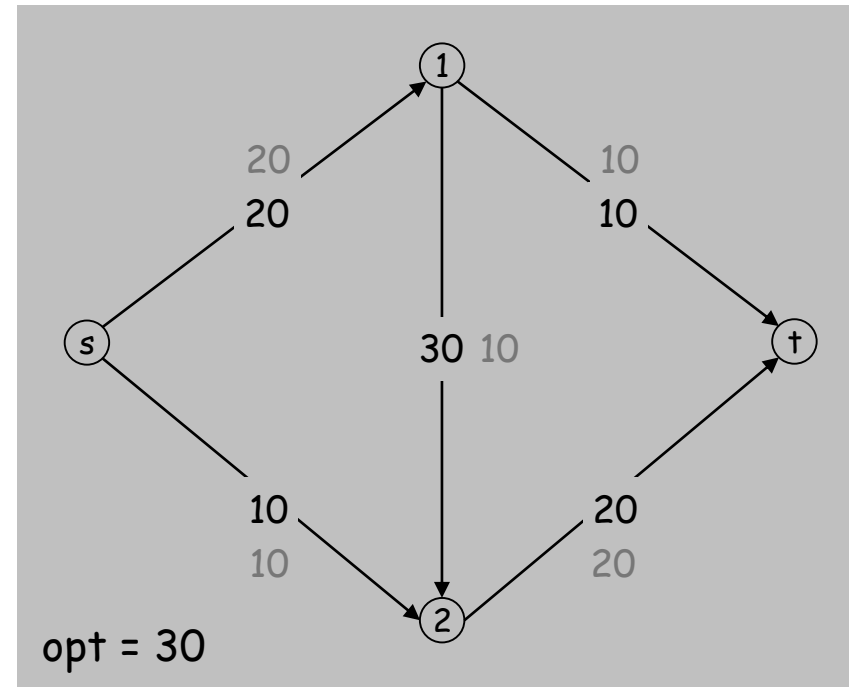
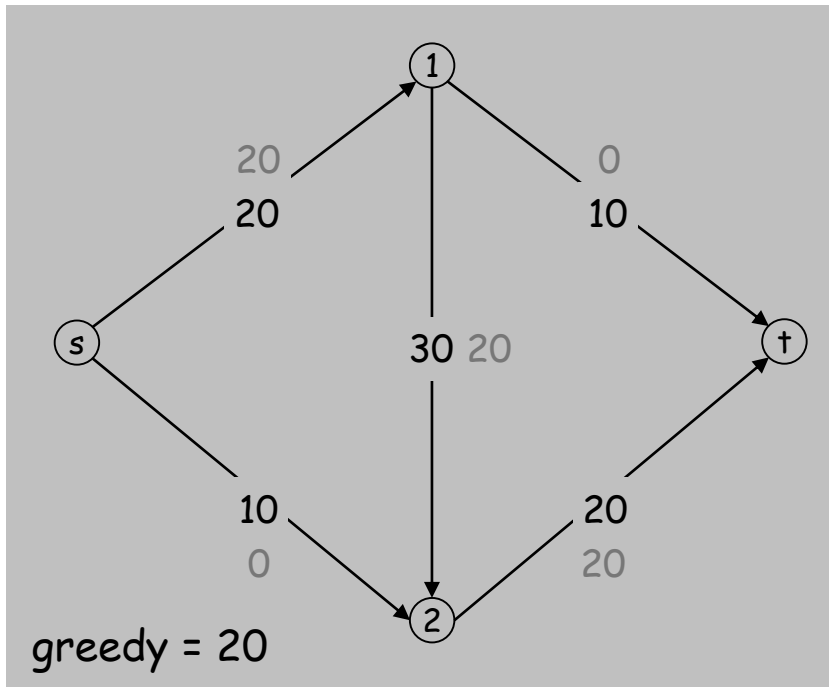
Flow value = 20

# Towards a Max Flow Algorithm

## Greedy algorithm.

- Start with  $f(e) = 0$  for all edge  $e \in E$ .
- Find an  $s$ - $t$  path  $P$  where each edge has  $f(e) < c(e)$ .
- Augment flow along path  $P$ .
- Repeat until you get **stuck**.

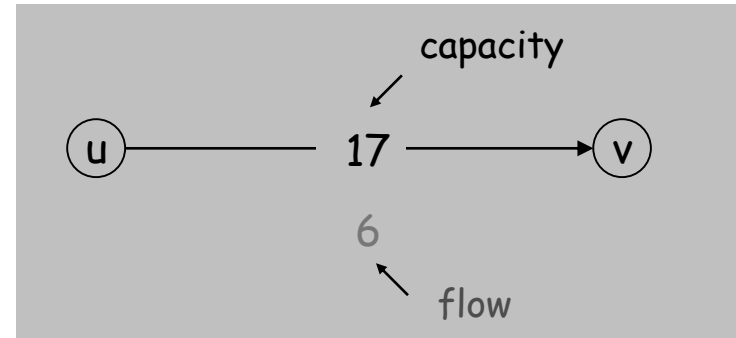
← locally optimality  $\nRightarrow$  global optimality



# Residual Graph

Original edge:  $e = (u, v) \in E$ .

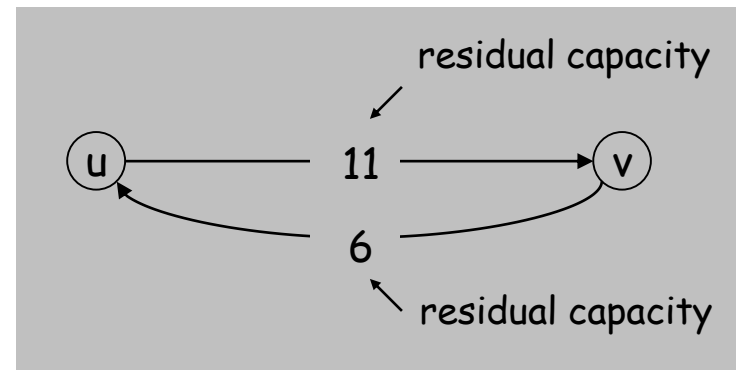
- Flow  $f(e)$ , capacity  $c(e)$ .



Residual edge.

- "Undo" flow sent.
- $e = (u, v)$  and  $e^R = (v, u)$ .
- Residual capacity:

$$c_f(e) = \begin{cases} c(e) - f(e) & \text{if } e \in E \\ f(e) & \text{if } e^R \in E \end{cases}$$

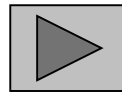
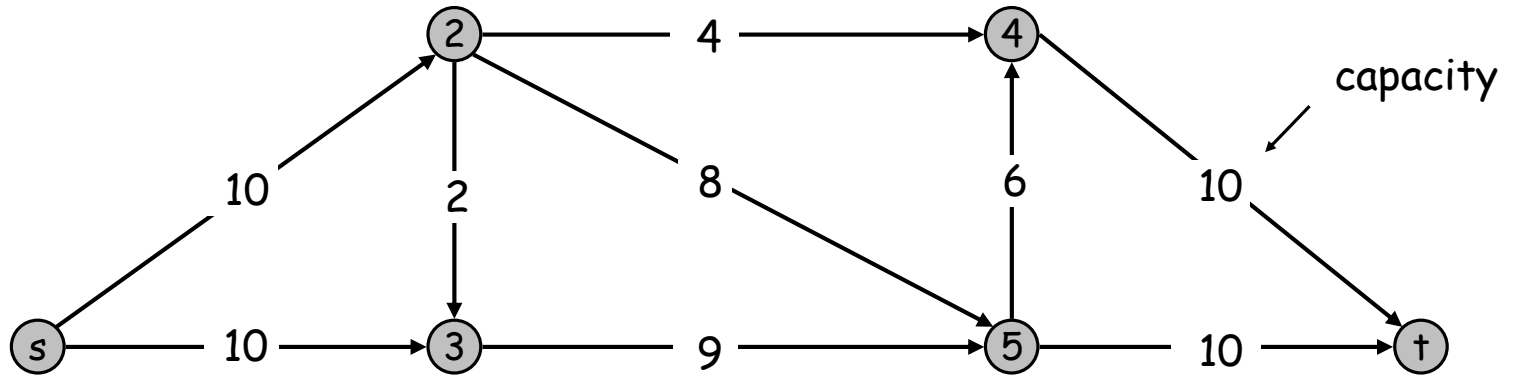


Residual graph:  $G_f = (V, E_f)$ .

- Residual edges with positive residual capacity.
- $E_f = \{e : f(e) < c(e)\} \cup \{e^R : f(e) > 0\}$ .

# Ford-Fulkerson Algorithm

$G$ :



# Augmenting Path Algorithm

```
Augment(f, c, P) {  
    b ← bottleneck(P)  
    foreach e ∈ P {  
        if (e ∈ E) f(e) ← f(e) + b  
        else      f(eR) ← f(eR) - b  
    }  
    return f  
}
```

forward edge  
reverse edge

```
Ford-Fulkerson(G, s, t, c) {  
    foreach e ∈ E f(e) ← 0  
    Gf ← residual graph  
  
    while (there exists augmenting path P) {  
        f ← Augment(f, c, P)  
        update Gf  
    }  
    return f  
}
```

# Max-Flow Min-Cut Theorem

**Augmenting path theorem.** Flow  $f$  is a max flow iff there are no augmenting paths.

**Max-flow min-cut theorem.** [Elias-Feinstein-Shannon 1956, Ford-Fulkerson 1956]  
The value of the max flow is equal to the value of the min cut.

**Pf.** We prove both simultaneously by showing TFAE:

- (i) There exists a cut  $(A, B)$  such that  $v(f) = \text{cap}(A, B)$ .
- (ii) Flow  $f$  is a max flow.
- (iii) There is no augmenting path relative to  $f$ .

(i)  $\Rightarrow$  (ii) This was the corollary to weak duality lemma.

(ii)  $\Rightarrow$  (iii) We show contrapositive.

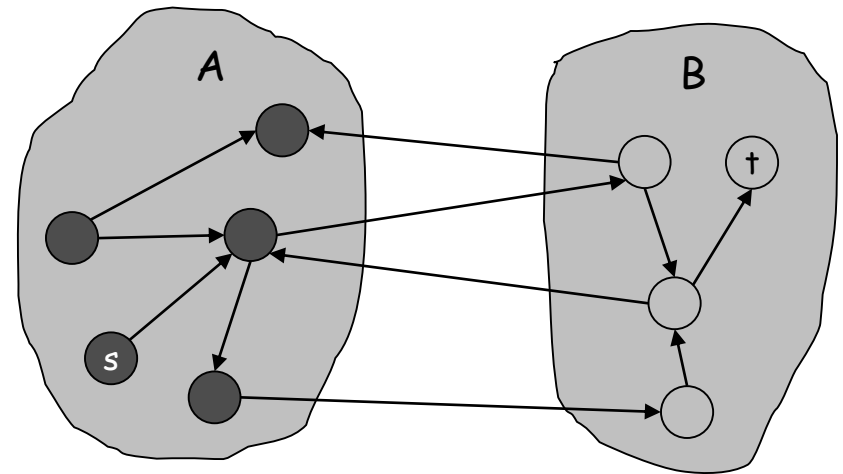
- Let  $f$  be a flow. If there exists an augmenting path, then we can improve  $f$  by sending flow along path.

# Proof of Max-Flow Min-Cut Theorem

(iii)  $\Rightarrow$  (i)

- Let  $f$  be a flow with no augmenting paths.
- Let  $A$  be set of vertices reachable from  $s$  in residual graph.
- By definition of  $A$ ,  $s \in A$ .
- By definition of  $f$ ,  $t \notin A$ .

$$\begin{aligned}v(f) &= \sum_{e \text{ out of } A} f(e) - \sum_{e \text{ in to } A} f(e) \\ &= \sum_{e \text{ out of } A} c(e) \\ &= \text{cap}(A, B) \quad \blacksquare\end{aligned}$$



original network

# Running Time

**Assumption.** All capacities are integers between 1 and  $C$ .

**Invariant.** Every flow value  $f(e)$  and every residual capacity  $c_f(e)$  remains an integer throughout the algorithm.

**Theorem.** The algorithm terminates in at most  $v(f^*) \leq nC$  iterations.

**Pf.** Each augmentation increase value by at least 1. ▀

**Corollary.** If  $C = 1$ , Ford-Fulkerson runs in  $O(mn)$  time.

**Integrality theorem.** If all capacities are integers, then there exists a max flow  $f$  for which every flow value  $f(e)$  is an integer.

**Pf.** Since algorithm terminates, theorem follows from invariant. ▀

# Choosing Good Augmenting Paths

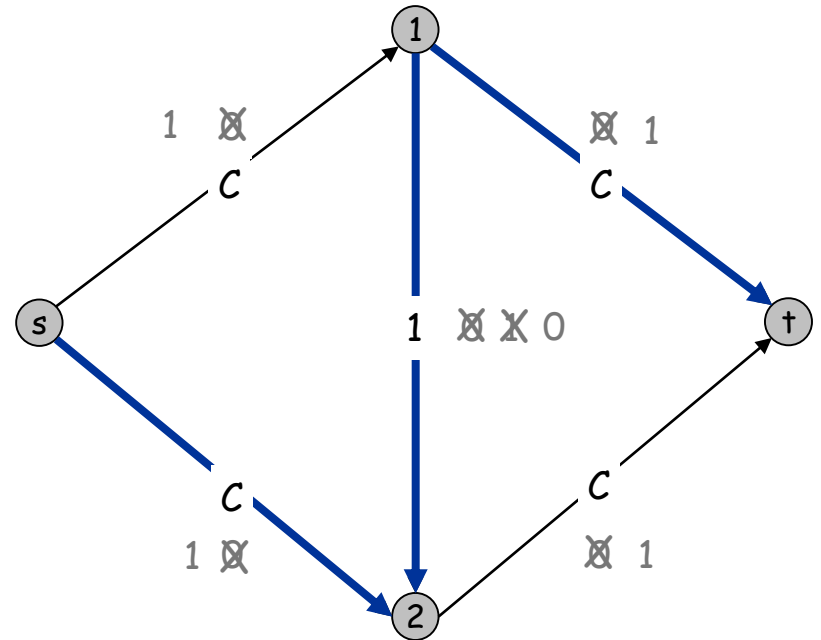
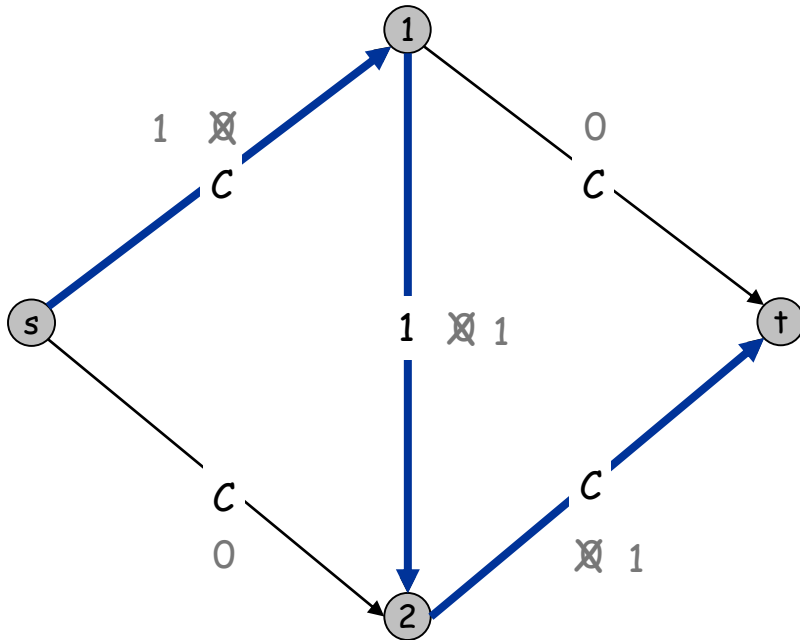
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# Ford-Fulkerson: Exponential Number of Augmentations

Q. Is generic Ford-Fulkerson algorithm polynomial in input size?

$m, n,$  and  $\log C$  ↗

A. No. If max capacity is  $C$ , then algorithm can take  $C$  iterations.



# Choosing Good Augmenting Paths

Use care when selecting augmenting paths.

- Some choices lead to exponential algorithms.
- Clever choices lead to polynomial algorithms.
- If capacities are irrational, algorithm not guaranteed to terminate!

Goal: choose augmenting paths so that:

- Can find augmenting paths efficiently.
- Few iterations.

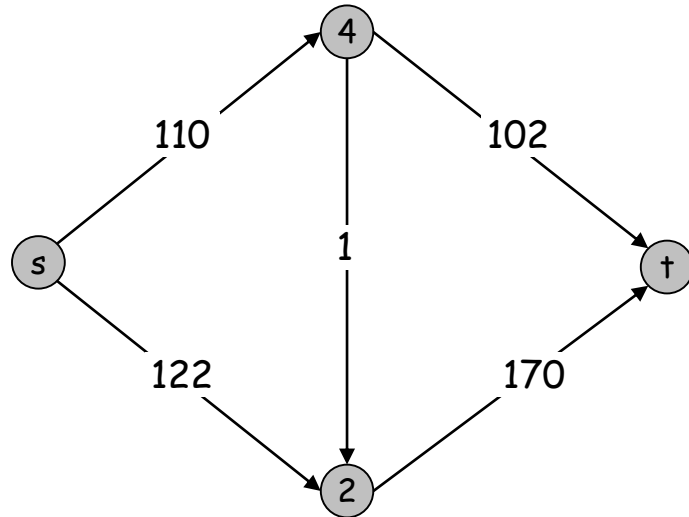
Choose augmenting paths with: [Edmonds-Karp 1972, Dinitz 1970]

- Max bottleneck capacity.
- Sufficiently large bottleneck capacity.
- Fewest number of edges.

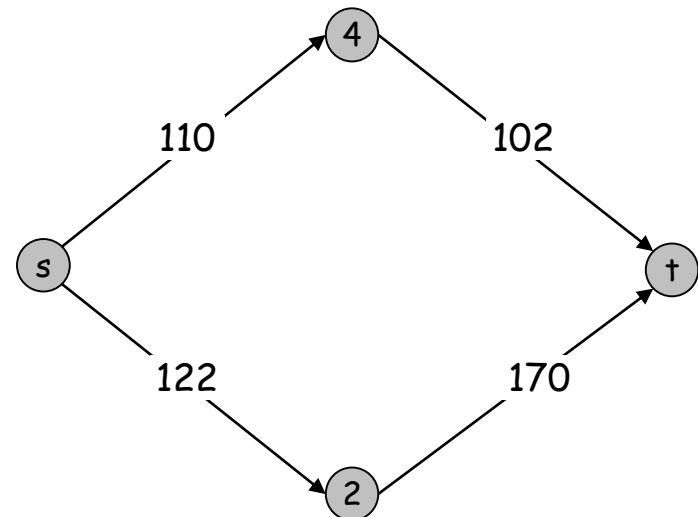
# Capacity Scaling

**Intuition.** Choosing path with highest bottleneck capacity increases flow by max possible amount.

- Don't worry about finding exact highest bottleneck path.
- Maintain scaling parameter  $\Delta$ .
- Let  $G_f(\Delta)$  be the subgraph of the residual graph consisting of only arcs with capacity at least  $\Delta$ .



$G_f$



$G_f(100)$

# Capacity Scaling

```
Scaling-Max-Flow( $G, s, t, c$ ) {  
  foreach  $e \in E$   $f(e) \leftarrow 0$   
   $\Delta \leftarrow$  smallest power of 2 greater than or equal to  $C$   
   $G_f \leftarrow$  residual graph  
  
  while ( $\Delta \geq 1$ ) {  
     $G_f(\Delta) \leftarrow \Delta$ -residual graph  
    while (there exists augmenting path  $P$  in  $G_f(\Delta)$ ) {  
       $f \leftarrow$  augment( $f, c, P$ )  
      update  $G_f(\Delta)$   
    }  
     $\Delta \leftarrow \Delta / 2$   
  }  
  return  $f$   
}
```

## Capacity Scaling: Correctness

**Assumption.** All edge capacities are integers between 1 and  $C$ .

**Integrality invariant.** All flow and residual capacity values are integral.

**Correctness.** If the algorithm terminates, then  $f$  is a max flow.

**Pf.**

- By integrality invariant, when  $\Delta = 1 \Rightarrow G_f(\Delta) = G_f$ .
- Upon termination of  $\Delta = 1$  phase, there are no augmenting paths. ▪

## Capacity Scaling: Running Time

**Lemma 1.** The outer while loop repeats  $1 + \lceil \log_2 C \rceil$  times.

**Pf.** Initially  $C \leq \Delta < 2C$ .  $\Delta$  decreases by a factor of 2 each iteration. ▀

**Lemma 2.** Let  $f$  be the flow at the end of a  $\Delta$ -scaling phase. Then the value of the maximum flow is at most  $v(f) + m \Delta$ . ← proof on next slide

**Lemma 3.** There are at most  $2m$  augmentations per scaling phase.

- Let  $f$  be the flow at the end of the previous scaling phase.
- $L2 \Rightarrow v(f^*) \leq v(f) + m (2\Delta)$ .
- Each augmentation in a  $\Delta$ -phase increases  $v(f)$  by at least  $\Delta$ . ▀

**Theorem.** The scaling max-flow algorithm finds a max flow in  $O(m \log C)$  augmentations. It can be implemented to run in  $O(m^2 \log C)$  time. ▀

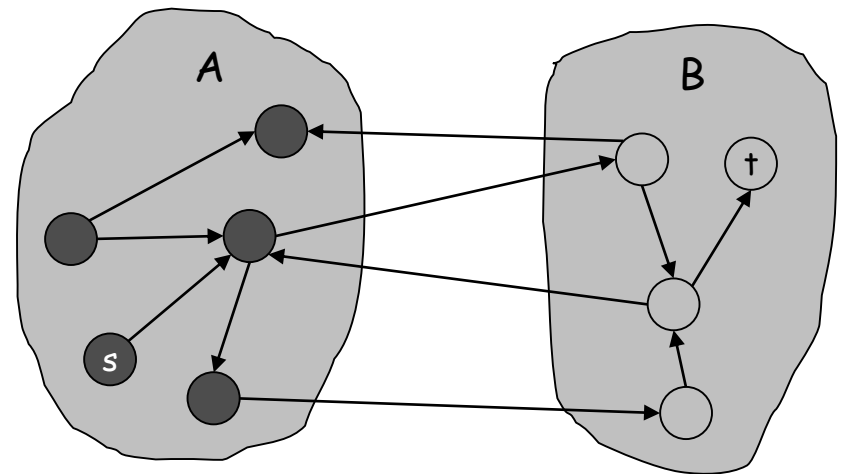
# Capacity Scaling: Running Time

**Lemma 2.** Let  $f$  be the flow at the end of a  $\Delta$ -scaling phase. Then value of the maximum flow is at most  $v(f) + m \Delta$ .

**Pf.** (almost identical to proof of max-flow min-cut theorem)

- We show that at the end of a  $\Delta$ -phase, there exists a cut  $(A, B)$  such that  $\text{cap}(A, B) \leq v(f) + m \Delta$ .
- Choose  $A$  to be the set of nodes reachable from  $s$  in  $G_f(\Delta)$ .
- By definition of  $A$ ,  $s \in A$ .
- By definition of  $f$ ,  $t \notin A$ .

$$\begin{aligned}
 v(f) &= \sum_{e \text{ out of } A} f(e) - \sum_{e \text{ in to } A} f(e) \\
 &\geq \sum_{e \text{ out of } A} (c(e) - \Delta) - \sum_{e \text{ in to } A} \Delta \\
 &= \sum_{e \text{ out of } A} c(e) - \sum_{e \text{ out of } A} \Delta - \sum_{e \text{ in to } A} \Delta \\
 &\geq \text{cap}(A, B) - m\Delta \quad \blacksquare
 \end{aligned}$$



original network

# References

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## References

- Sections 7.1-3 of the text book "algorithm design" by Jon Kleinberg and Eva Tardos
- The original slides were prepared by Kevin Wayne. The slides are distributed by Pearson Addison-Wesley.